

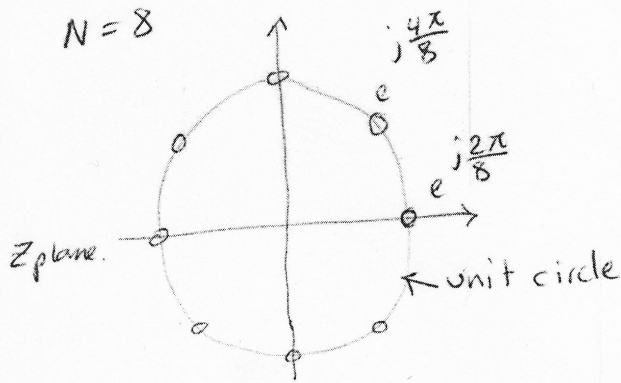
**problem 8.1**

Show that each of the numbers  $e^{j(\frac{2\pi}{N})k}$ ,  $0 \leq k \leq N-1$  corresponds to an  $N$ th root of unity. Plot these numbers as phasors in the complex plane and illustrate, by means of this figure, the orthogonality property.

$$\sum_{n=0}^{N-1} e^{j(\frac{2\pi}{N})kn} e^{-j(\frac{2\pi}{N})ln} = \begin{cases} N, & \text{if } \text{mod}(k-l, N) = 0 \\ 0, & \text{otherwise} \end{cases}$$

let  $\left[ e^{j(\frac{2\pi}{N})k} \right]^N = e^{j2\pi k} = 1 \Rightarrow X^N = 1$  i.e.  $e^{j\frac{2\pi k}{N}}$  is the

$N$ th root of unity.



⊕ now consider:

$$\sum_{n=0}^{N-1} e^{j(\frac{2\pi}{N})kn} e^{-j(\frac{2\pi}{N})ln}$$

→ if  $k \neq l$ , the sum represents the  $N$  equally spaced roots on the unit circle, the roots sum to zero.

→ if  $k = l$ ,

$$\text{sum} \sum_{n=0}^{N-1} e^{j\frac{2\pi}{N}(k-l)n} = N$$

**problem 8.3**

let  $x(n)$  be a real-valued  $N$ -point ( $N=2^v$ ) sequence. Develop a method to compute an  $N$ -point DFT  $X'(k)$ , which contains only the odd harmonics [i.e.  $X'(k) = 0$  if  $k$  is even] by using only a real  $N/2$ -point DFT.

$$\begin{aligned} X(k) &= \sum_{n=0}^{N-1} x(n) W_N^{kn}, \quad 0 \leq k \leq N-1 \\ &= \sum_{n=0}^{\frac{N}{2}-1} x(n) W_N^{kn} + \sum_{n=\frac{N}{2}}^{N-1} x(n) W_N^{kn} \\ &= \sum_{n=0}^{\frac{N}{2}-1} x(n) W_N^{kn} + \sum_{r=0}^{\frac{N}{2}-1} x(r+\frac{N}{2}) W_N^{(r+\frac{N}{2})k} \end{aligned}$$

so to keep odd harmonics,

$$\text{let } X'(k') = X(2k'), \quad 0 \leq k' \leq \frac{N}{2}-1$$

$$\Rightarrow X'(k') = \sum_{n=0}^{\frac{N}{2}-1} x(n) W_N^{(2k'+1)n} + \sum_{r=0}^{\frac{N}{2}-1} x(r+\frac{N}{2}) W_N^{(r+\frac{N}{2})(2k'+1)}$$

knowing  $W_N^{2k'n} = W_{\frac{N}{2}}^{k'n}$ ,  $W_N^N = 1$

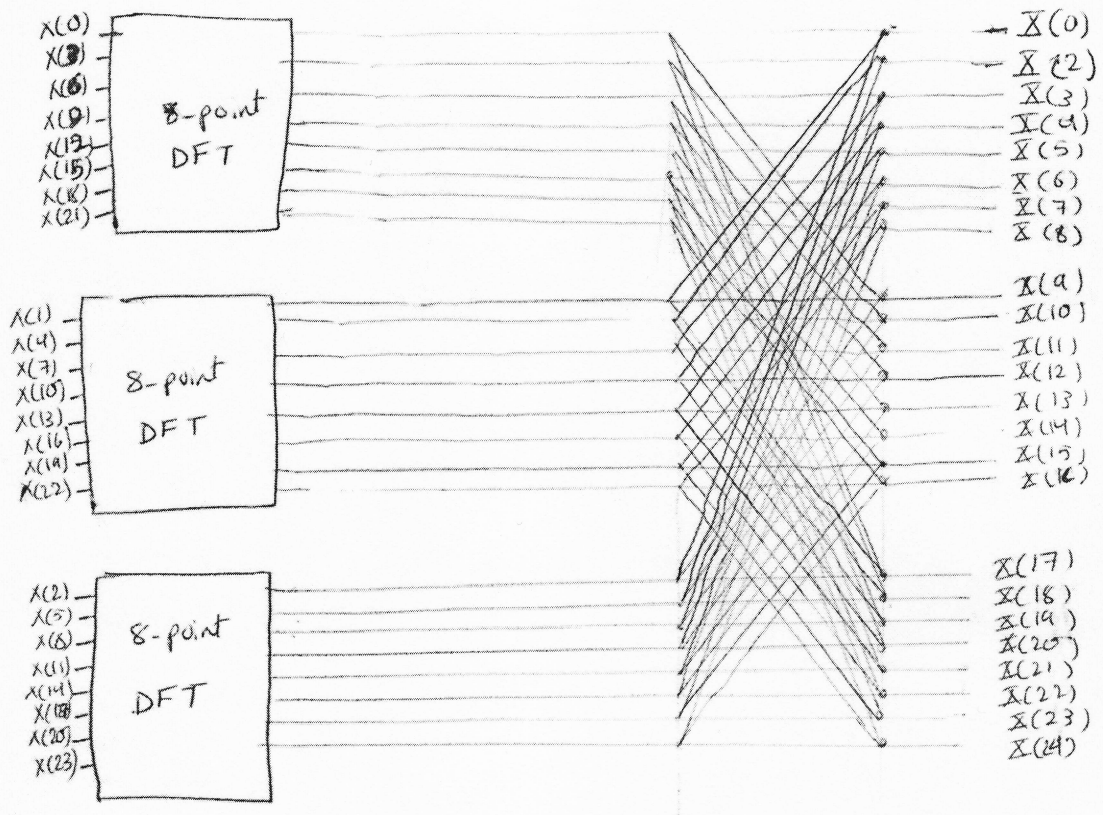
$$\Rightarrow X'(k') = \sum_{n=0}^{\frac{N}{2}-1} \left[ x(n) W_N^{k'n} W_{\frac{N}{2}}^{k'n} + x(n+\frac{N}{2}) W_N^{k'n} W_{\frac{N}{2}}^{k'n} W_N^{N \cdot \frac{k'n}{2}} \right]$$

$W_N^{N \cdot \frac{k'n}{2}} = e^{-j \frac{2\pi}{N} \cdot \frac{N}{2} \cdot \frac{k'n}{2}} = -1$

$$X'(k') = \sum_{n=0}^{\frac{N}{2}-1} \left[ x(n) - x(n+\frac{N}{2}) \right] W_N^{k'n} W_{\frac{N}{2}}^{k'n}$$

Problem 8.4 A designer has available a number of eight-point FFT chips. Show explicitly how he would interconnect three such chips in order to compute the 24-point DFT

$$\begin{aligned}
 \bar{X}(k) &= \sum_{n=0,3,6,\dots}^{21} x(n) W_N^{kn} + \sum_{n=1,4,7,\dots}^{22} x(n) W_N^{kn} + \sum_{n=2,5,8,\dots}^{23} x(n) W_N^{kn} \\
 &= \sum_{r=0}^7 x(3r) W_{\frac{N}{3}}^{kr} + \sum_{r=0}^7 x(3r+1) W_N^{(3r+1)k} + \sum_{r=0}^7 x(3r+2) W_N^{(3r+2)k} \\
 &= \sum_{r=0}^7 x(3r) W_{\frac{N}{3}}^{kr} + \sum_{r=0}^7 x(3r+1) W_{\frac{N}{3}}^{kr} W_N^k + \sum_{r=0}^7 x(3r+2) W_{\frac{N}{3}}^{kr} W_N^{2k} \\
 &\equiv F_1(k) + F_2(k) W_N^k + F_3(k) W_N^{2k}
 \end{aligned}$$



$$e^{-j\frac{2\pi}{N}k}$$



**problem 8.8** Compute the eight-point DFT of the sequence

$$x(n) = \begin{cases} 1, & 0 \leq n \leq 7 \\ 0, & \text{otherwise} \end{cases}$$

by using the decimation-in-frequency FFT algorithm described in text.

following fig. 8.1.11.

1<sup>st</sup> stage  $\rightarrow \{2, 2, 2, 2, 0, 0, 0, 0\}$

2<sup>nd</sup> stage  $\rightarrow \{4, 4, 0, 0, 0, 0, 0, 0\}$

3<sup>rd</sup> stage  $\rightarrow \{8, 0, 0, 0, 0, 0, 0, 0\}$

reordering  $\rightarrow$  remains  $\{8, 0, 0, 0, 0, 0, 0, 0\}$



problem 8.81 Compute the 8-point DFT of the sequence

$$x(n) = \left\{ \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, 0, 0, 0, 0 \right\}$$

using the in-place 2 radix-2 decimation-in-time & radix-2 decimation in-frequency algorithms. Follow exactly the corresponding signal flow graphs and keep track of all the intermediate quantities by putting them on the diagrams.

Figure 8.16 Decimation-in-time FFT

$$\rightarrow 1^{\text{st}} \text{ stage output} \rightarrow \left\{ \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \dots, \frac{1}{2} \right\}$$

$$\rightarrow 2^{\text{nd}} \text{ stage output} \rightarrow \left\{ \underbrace{\left( \frac{1}{2} + \frac{1}{2} \cdot w_8^0 \right)}, \frac{1}{2} (1 + w_8^2), \frac{\frac{1}{2} - \frac{1}{2} w_8^0}{0}, \frac{1}{2} (1 - w_8^2), \right. \\ \left. \frac{1}{2} (1 + w_8^2), 0, \frac{1}{2} (1 - w_8^2) \right\}$$

$$\rightarrow 3^{\text{rd}} \text{ stage output} \rightarrow \left\{ 2, 0, 0, 0, \frac{1}{2} (1 + w_8^1 + w_8^2 + w_8^3), \frac{1}{2} (1 - w_8^1 + w_8^2 - w_8^3), \right. \\ \left. \frac{1}{2} (1 - w_8^2 + w_8^3 - w_8^5), \frac{1}{2} (1 - w_8^2 - w_8^3 + w_8^5) \right\}$$

following Figure 8.1.11 for Decimation-in-frequency FFT.

$$\rightarrow 1^{\text{st}} \text{ stage: } \left\{ \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \frac{1}{2}, \frac{1}{2} w_8^1, \frac{1}{2} w_8^2, \frac{1}{2} w_8^3 \right\}$$

$$\rightarrow 2^{\text{nd}} \text{ stage: } \left\{ 1, 1, 0, 0, \frac{1}{2} + \frac{1}{2} w_8^2, \frac{1}{2} w_8^1 + \frac{1}{2} w_8^3, \frac{1}{2} + \frac{1}{2} w_8^2, \frac{1}{2} w_8^3 - \frac{1}{2} w_8^5 \right\}$$

$$\rightarrow 3^{\text{rd}} \text{ stage: } \left\{ 2, 0, 0, 0, \frac{1}{2} (1 + w_8^1 + w_8^2 + w_8^3), \frac{1}{2} (1 - w_8^1 + w_8^2 - w_8^3), \right. \\ \left. \frac{1}{2} (1 - w_8^2 + w_8^3 - w_8^5), \frac{1}{2} (1 - w_8^2 - w_8^3 + w_8^5) \right\}$$

**8.13** Consider the eight-point decimation-in-time algorithm (DIT) flow graph in 8.1.6.

(a) What is the gain of the "signal path" that goes from  $X(1)$  to  $X(2)$ ?

$$\Rightarrow \text{"gain"} = W_8^0 W_8^0 (-1) W_8^2 = -W_8^2 = -e^{-j\frac{2\pi \cdot 2}{8}} = j$$

(b) How many paths lead from the input to a given output sample? Is this true for every output sample?

$\Rightarrow$  Given a certain output sample, there is one path from every input leading to it. This is true for every output.

(c) Compute  $X(3)$  using the operations dictated by this flow graph.

$$\Rightarrow X(3) = X(0) + W_8^3 X(1) - W_8^2 X(2) + W_8^2 W_8^3 X(3) - W_8^0 X(4) - W_8^0 W_8^3 X(5) + W_8^0 W_8^2 X(6) + W_8^0 W_8^2 W_8^3 X(7)$$

**8.16** Show that the product of two complex numbers  $(a+jb)$  and  $(c+jd)$  can be performed with three real ~~number~~ multiplications and five additions using the algorithm

$$X_R = (a-b)d + (c-d)a$$

$$X_I = (a-b)d + (c+d)b$$

$$\text{Where } X = X_R + jX_I = (a+jb)(c+jd)$$

$\Rightarrow$  From the algorithm,  $(a-b)d$  is common for computation of  $X_R$  and  $X_I$ . We have  $(a-b)d$ ,  $(c-d)a$ , and  $(c+d)b$ , each of which has 1 add and 1 multiplication. We need one more add for  $X_R$  and  $X_I$ , respectively.

$$\Rightarrow 3 \times 1 + 2 = 5 \text{ add and } 3 \times 1 = 3 \text{ multiplication.}$$

**8.19** Let  $X(k)$  be the  $N$ -point DFT of the sequence  $x(n)$ ,  $0 \leq n \leq N-1$ . What is the  $N$ -point DFT of the sequence  $s(n) = X(n)$ ,  $0 \leq n \leq N-1$ ?

$\Rightarrow$  Given  $X(k) = \sum_{n=0}^{N-1} x(n) W_N^{kn}$ . Let  $F(t)$  be the DFT of the sequence  $s(n) = X(n)$ .

$$F(t) = \sum_{k=0}^{N-1} X(k) W_N^{tk} = \sum_{k=0}^{N-1} \left( \sum_{n=0}^{N-1} x(n) W_N^{kn} \right) W_N^{tk} = \sum_{n=0}^{N-1} x(n) \left[ \sum_{k=0}^{N-1} W_N^{kn} W_N^{tk} \right]$$

$$= \sum_{n=0}^{N-1} x(n) \delta[(n+t)]_N = \sum_{n=0}^{N-1} x(n) \delta[N-1-n-t] \quad t=0, 1, \dots, N-1$$

$$= \{x(N-1), x(N-2), \dots, x(1), x(0)\}$$

8.20 Let  $X(k)$  be the  $N$ -point DFT of the sequence  $x(n)$ ,  $0 \leq n \leq N-1$ . We define a  $2N$ -point sequence  $y(n)$  as  $y(n) = \begin{cases} x(\frac{n}{2}), & n \text{ even} \\ 0, & n \text{ odd} \end{cases}$ . Express the  $2N$ -point DFT of  $y(n)$  in terms of  $X(k)$ .

$$\begin{aligned} \Rightarrow Y(k) &= \sum_{n=0}^{2N-1} y(n) W_{2N}^{kn} = \sum_{n=0, n \text{ even}}^{2N-1} y(n) W_{2N}^{kn} = \sum_{m=0}^{N-1} y(2m) W_{2N}^{2km} \\ &= \sum_{m=0}^{N-1} y(2m) W_N^{km} = \sum_{m=0}^{N-1} x(m) W_N^{km} = X(k), \quad k \in [0, N-1] \\ &= X(k-N), \quad k \in [N, 2N-1]. \end{aligned}$$

8.35 The basic butterfly in the radix-2 decimation-in-time FFT algorithm is

$$X_{n+1}(k) = X_n(k) + W_N^m X_n(l)$$

$$X_{n+1}(l) = X_n(k) - W_N^m X_n(l)$$

(a) If we require that  $|X_n(k)| < \frac{1}{2}$ , and  $|X_n(l)| < \frac{1}{2}$ , show that

$$|\operatorname{Re}[X_{n+1}(k)]| < 1, \quad |\operatorname{Re}[X_{n+1}(l)]| < 1$$

$$|\operatorname{Im}[X_{n+1}(k)]| < 1, \quad |\operatorname{Im}[X_{n+1}(l)]| < 1.$$

$\Rightarrow$  This overflow does not occur.

$$\Rightarrow \operatorname{Re}[X_{n+1}(k)] = \frac{1}{2} X_{n+1}(k) + \frac{1}{2} X_{n+1}^*(k)$$

$$= \frac{1}{2} X_n(k) + \frac{1}{2} W_N^m X_n(l) + \frac{1}{2} X_n^*(k) - \frac{1}{2} W_N^{-m} X_n^*(l)$$

$$= \operatorname{Re}[X_n(k)] + \operatorname{Re}[W_N^m X_n(l)]$$

$$\text{Since } \begin{cases} |X_n(k)| < \frac{1}{2} \\ |X_n(l)| < \frac{1}{2} \end{cases} \Rightarrow \begin{cases} |\operatorname{Re}[X_n(k)]| < \frac{1}{2} \\ |\operatorname{Re}[X_n(l) W_N^m]| < \frac{1}{2} \end{cases}$$

$$\text{So } |\operatorname{Re}[X_{n+1}(k)]| \leq |\operatorname{Re}[X_n(k)]| + |\operatorname{Re}[W_N^m X_n(l)]| < 1$$

The other inequalities are verified similarly.

(b) Prove that  $\max[|X_{n+1}(k)|, |X_{n+1}(l)|] \geq \max[|X_n(k)|, |X_n(l)|]$ .

$$\max[|X_{n+1}(k)|, |X_{n+1}(l)|] \leq 2 \max[|X_n(k)|, |X_n(l)|]$$

$$\Rightarrow X_{n+1}(k) = \left\{ \operatorname{Re}[X_n(k)] + j \operatorname{Im}[X_n(k)] \right\} + \left\{ \cos\left(\frac{2\pi}{N}m\right) - j \sin\left(\frac{2\pi}{N}m\right) \right\} \left\{ \operatorname{Re}[X_n(l)] + j \operatorname{Im}[X_n(l)] \right\}$$



$$= \text{Re}[X_n(k)] + \cos\left(\frac{2\pi}{N}m\right) \text{Re}[X_n(l)] + \sin\left(\frac{2\pi}{N}m\right) \text{Im}[X_n(l)] + j \{ \text{Im}[X_n(k)] + \cos\left(\frac{2\pi}{N}m\right) \text{Im}[X_n(l)] + \sin\left(\frac{2\pi}{N}m\right) \text{Re}[X_n(l)] \}$$

Therefore,  $|X_{n+1}(k)| = |X_n(k)| + |X_n(l)| + A,$

where  $A = 2 \cos\left(\frac{2\pi}{N}m\right) \{ \text{Re}[X_n(k)] \text{Re}[X_n(l)] + \text{Im}[X_n(k)] \text{Im}[X_n(l)] \} + 2 \sin\left(\frac{2\pi}{N}m\right) \{ \text{Re}[X_n(k)] \text{Im}[X_n(l)] - \text{Im}[X_n(k)] \text{Re}[X_n(l)] \}$

also  $|X_{n+1}(l)|^2 = |X_n(k)|^2 + |X_n(l)|^2 - A$

Therefore, if  $A \geq 0$

$$\begin{aligned} \max [ |X_{n+1}(k)|, |X_{n+1}(l)| ] &= |X_{n+1}(k)| \\ &= \{ |X_n(k)|^2 + |X_n(l)|^2 + A \}^{\frac{1}{2}} \\ &\geq \max [ |X_n(k)|, |X_n(l)| ] \end{aligned}$$

We can prove that the inequality hold if  $A < 0$ .

Also, from the pair of equation ~~for~~ for computing the butterfly outputs, we have

$$2X_n(k) = X_{n+1}(k) + X_{n+1}(l)$$

$$2X_n(l) = W_N^{-m} X_{n+1}(k) - W_N^{-m} X_{n+1}(l).$$

Based above two equations, by a similar method to that employed above, it can be shown that

$$2 \max [ |X_n(k)|, |X_n(l)| ] \geq \max [ |X_{n+1}(k)|, |X_{n+1}(l)| ]$$

**problem 8.25** Develop a radix-3 decimation-in-time FFT algorithm for  $N=3^2$ , and draw the corresponding flow graph for  $N=9$ . What is the number of required complex multiplications? can the operations be performed in place?

$$\begin{aligned}
 X(k) &= \sum_{n=0}^8 x(n) W_9^{nk} = \sum_{n=0,3,6} x(n) W_9^{nk} + \sum_{n=1,4,7} x(n) W_9^{nk} + \sum_{n=2,5,8} x(n) W_9^{nk} \\
 &= \sum_{m=0}^{\frac{N}{3}-1=2} x(3m) W_9^{(3m)k} + \sum_{m=0}^2 x(3m+1) W_9^{(3m+1)k} + \sum_{m=0}^2 x(3m+2) W_9^{(3m+2)k} \\
 &= \sum_{m=0}^2 x(3m) W_{\frac{9}{3}}^{mk} + W_9^k \sum_{m=0}^2 x(3m+1) W_3^{km} + W_9^{2k} \sum_{m=0}^2 x(3m+2) W_3^{mk} \\
 &= F_1(k) + W_9^k F_2(k) + W_9^{2k} F_3(k), \quad 0 \leq k \leq N-1
 \end{aligned}$$

⊛ exploiting the periodicity of  $F_1(k), F_2(k), F_3(k)$ , for  $0 \leq k \leq \frac{N}{3}-1$

$$\begin{aligned}
 X(k) &= F_1(k) + W_9^k F_2(k) + W_9^{2k} F_3(k) \\
 X(k + \frac{N}{3}) &= F_1(k) + W_3^1 W_9^k F_2(k) + W_3^2 W_9^{2k} F_3(k) \\
 X(k + \frac{2N}{3}) &= F_1(k) + W_3^2 W_9^k F_2(k) + W_3^4 W_9^{2k} F_3(k)
 \end{aligned}$$

⊛ we have  $\log_3 N = 2$  stages.

# of complex multiplications:

⊛ in each stage  $\rightarrow$  6 complex multiplication by twiddle factors + 12 multiplications inside the butterfly's.

$\Rightarrow$  total # of complex multiplications = 36.

⊛ since  $W^0 = 1 \Rightarrow$  # of multiplications = 28

