

# Secure Information Aggregation for Smart Grids using Homomorphic Encryption

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## Outline

- Motivation, approach
- Related works
- Background:
  - ✦ Homomorphic Encryption
  - ✦ Honest-but-curious Model
- Secure Information Aggregation
- Example
- Analysis and discussion
- Personal insight and assessment
- References

## Motivation, approach

- Instant aggregation of power usage data:
  - ✦ At **different levels**: Neighborhood , subdivision, district, city etc. and at **different frequencies**.
- Essential for:
  - ✦ **Monitoring and predicting** power consumption.
  - ✦ **Allocating and balancing** loads and resources.
  - ✦ Administering power generation, etc.
- **Goal**: efficient and secure data aggregation for smart grids.
- Approach:
  - ✦ **In-network** distributed aggregation.
  - ✦ **Homomorphic** encryption.

## Related Works

- Various in-network data aggregation approaches:
  - ✦ For sensor networks, sensors are **limited by battery and resources**.
  - ✦ Sensors in the network are usually **trusted**, and security is against eavesdroppers and tampering attacks using fake inputs.
- Smart Grids:
  - ✦ Power of the smart meter is not a concern, but **communication bandwidth** is, specially when frequent aggregation is required.
  - ✦ Power usage is considered a **privacy** of the user.
  - ✦ **\*Traditional tree-based aggregation on plaintext does not apply.**

## Background: Homomorphic Encryption

### • Homomorphic encryption:

- ✦ A form of encryption where specific algebraic operation performed on the plaintext is **equivalent to another (possibly different)** algebraic operation performed on the ciphertext.
- ✦ Given a homomorphic encryption function  $E()$ , and two messages  $x, y \in Z_N$

$$E_k(x \star y) = E_{k1}(x) \circ E_{k2}(y)$$

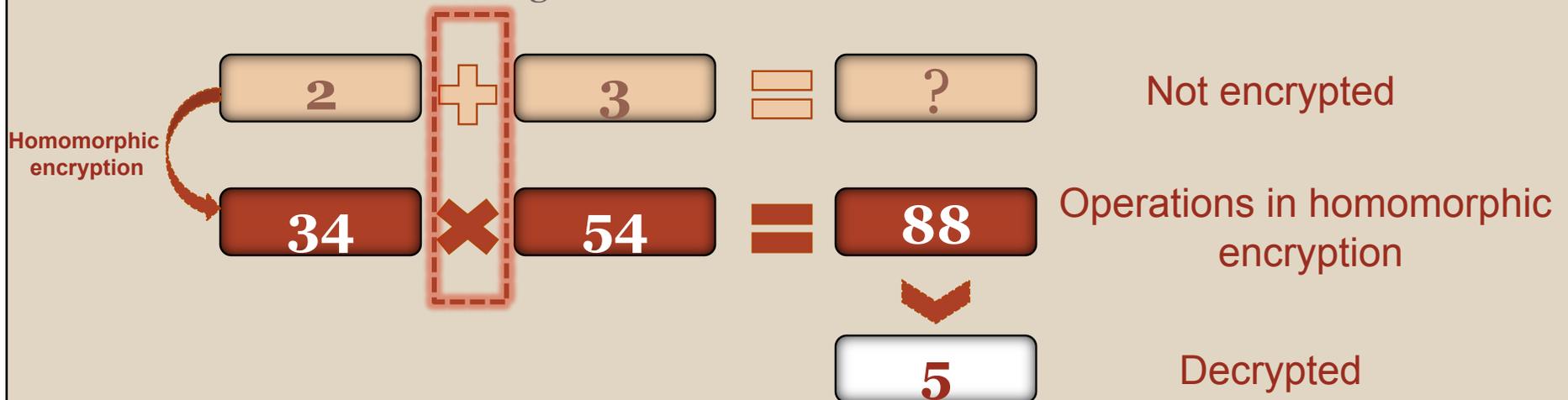
Without knowing the plaintext  $x, y$  or the private key.

- ✦ Used for **privacy-preserving operations**, voting.
- ✦ **Known schemes:** RSA, El Gamal, Paillier, Naccache-Stern, BGN etc.
- ✦ Paper **adopts Paillier** scheme.

# Homomorphic encryption

- Paillier cryptosystem:

- ✦ Invented in 1999 by Pascal Pailier.
- ✦ Has additive homomorphic property.
  - Given only the public-key and the encryption of  $m_1$  and  $m_2$ , one can compute the encryption of  $m_1 + m_2$ . [2]
- ✦ Indeterministic:
  - the same message will be encrypted into different ciphers using different random blinding factors.

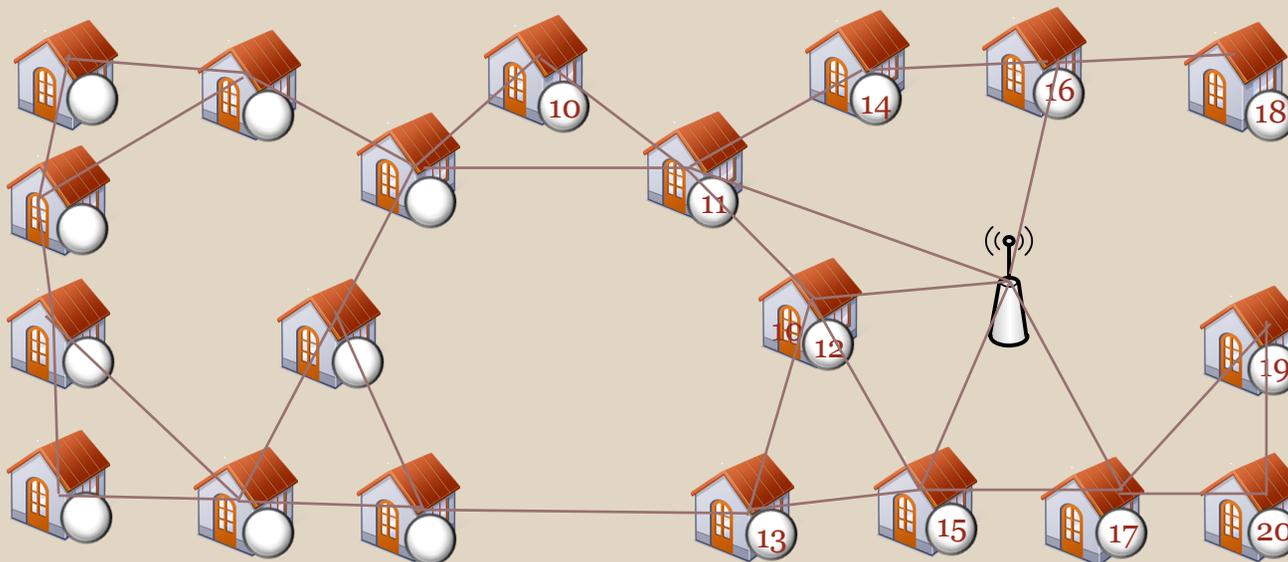


## Background: Honest-but-curious Model

- Honest-but-curious model:
  - ✦ All parties are assumed to follow protocol properly “**honest**”.
  - ✦ Keep all inputs from other parties and all intermediate computation results “**curious**”.
- Honest-but-curious **smart meters**:
  - ✦ Do not tamper with the aggregation protocols
  - ✦ Do not drop or distort any source value or intermediate result.
  - ✦ Will try to infer others’ electricity usage from messages routed through them

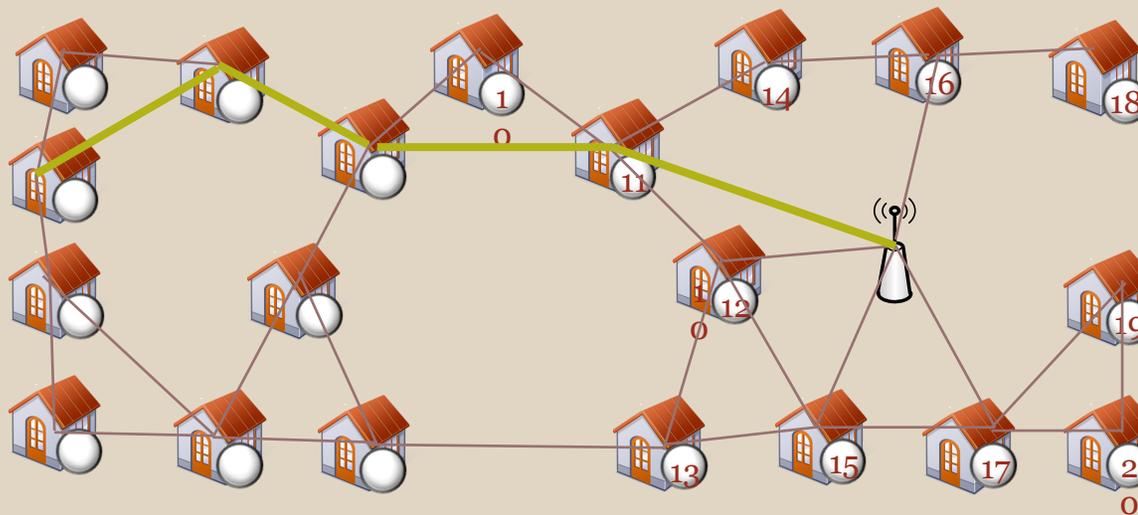
# Secure Information Aggregation

- Smart Grid Communication Infrastructure:
  - ✦ Most popular: **wireless-wired** multi-layer architecture.
    - Wireless: smart meters in a neighborhood communication with a collector device.
    - Wired: collector device with the rest of the grid.



## Secure Information Aggregation

- Data Aggregation: important type of query in Smart Grids.
  - ✦ Example: average power usage of the neighborhood.
  - ✦ Traditionally: every smart meter establishes a connection with the collector and uses it exclusively to report its data.
    - Excessive network traffic.
    - Overwhelming demands at the collectors.



# Secure in-network incremental aggregation

- Approach:
  - ✦ Establish an **aggregation tree**.
  - ✦ Enroute meters to **share** the channel.
  - ✦ **Ensure privacy** using homomorphic encryption.
    - With reasonable computation overhead.

# The Aggregation Tree

- To enable in-network Aggregation:
  - ✦ Aggregation path:
    - All smart meters in the neighborhood.
  - ✦ For each aggregation task:
    - All or subset of nodes on the aggregation path participate.

# The Aggregation Tree

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- Considering the smart meter network as a graph:
  - ✦ Graph  $G(V, E)$ :
    - $V$ , set of smart meters (vertices).
    - $E$ , set of available wireless links (edges) between any two smart meters.
    - Graph should be connected; every smart meter should have at least one communication path to the collector.

## The Aggregation Tree (cont.)

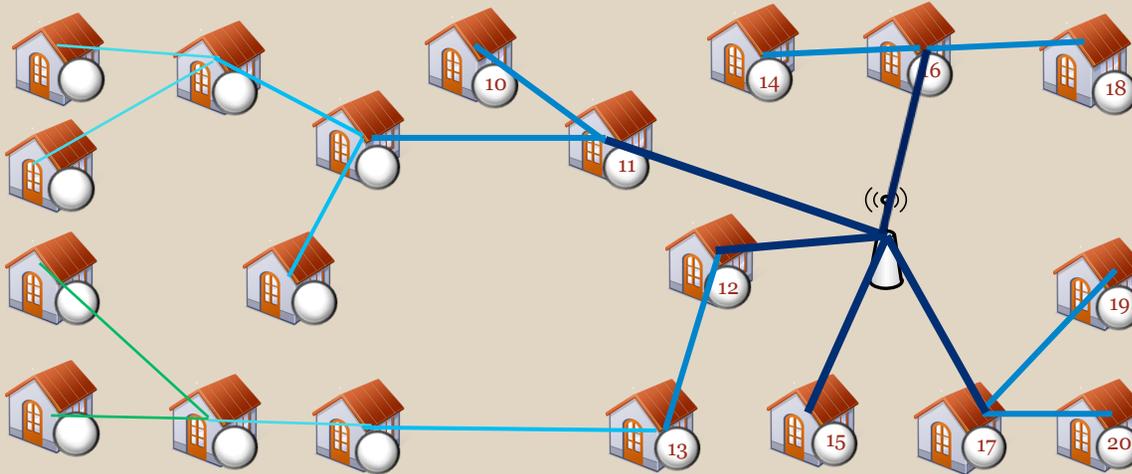
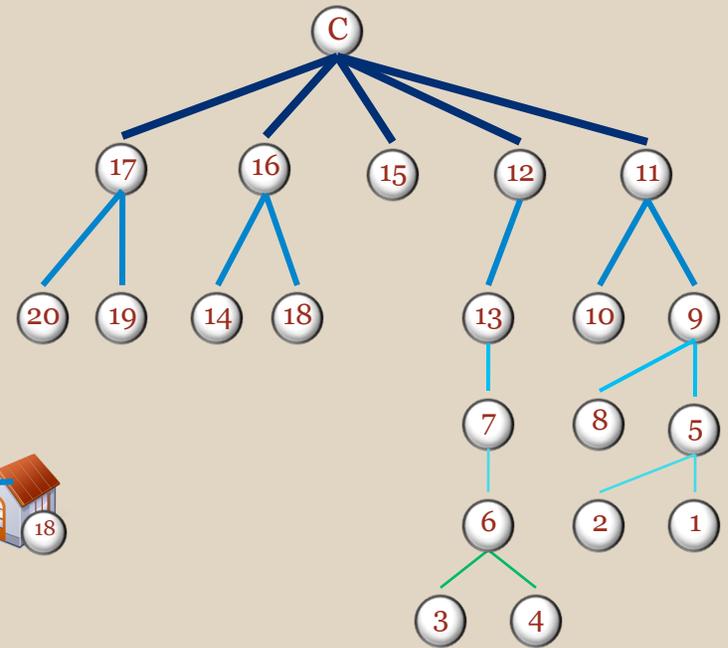
- - ✦ The Aggregation Tree:
    - A spanning tree of the graph with **minimal** subset of  $E$  that **connects all the** vertices.
    - Always **roots at the collector node**; which **initializes** all aggregation tasks and collects the results.
    - Aggregation is recursively calculated in a **bottom-up manner**; every nodes takes input from itself and its children nodes, aggregate the data and sends the result to its parent node.
  - ✦ Collector device:
    - Has the network graph of the entire neighborhood.
    - The aggregation tree is constructed locally at the collector.
    - An aggregation tree remains valid for an extended period of time.

# Constructing the Aggregation Tree

- ✦ Algorithm goals:
  - Height of the tree should be **small**.
  - An interior node should **not** have **too many** children, to avoid excessive computation and communication load.
- ✦ Approach:
  - **Breadth-first traversal** of the graph, starting at the collector node.
  - If node  $K$  has too many children **rebalance** the tree.
    - If a child of  $K$  is connected to a less populated sibling of  $K$ , move child to that sibling (will not increase the height of the tree).
    - If  $K$  still has too many children, check if a child is also connected to another child of  $K$ , and move it to that child (may increase the height of the tree).

# Example: constructing the Aggregation Tree

Aggregation tree constructed from the graph



Breadth-first traversal of the network graph

## In-network aggregation using homomorphic encryption

- Having the aggregation tree:
  - ✦ Construct operation plans for participating nodes (smart meters).
  - ✦ Deploy the operation plans in a top-down manner.

## In-network aggregation using homomorphic encryption

- An operation plan for a smart meter:

$\{T_{ID}, Trigger, Data, Collect, Operation, Destination, Key\}$

*T<sub>ID</sub>*, arbitrary unique identifier to identify message.

*Trigger*, defines when the aggregation will be conducted; periodically, upon collector request, or at a particular time. Time of local data reading, important in time-sensitive tasks.

*Data*, what information from the local smart grid will be collected in the aggregation.

*Collect*, tells a smart meter to wait for input from its children in the aggregation.

*Operation*, what operation to be performed; pre-processing, encryption and operations for aggregation.

*Destination*, the parent node, to whom the output from *Operation* will be submitted.

*Key*, a public key from the collector to be used to encrypt the local data.

## In-network aggregation using homomorphic encryption

- Output message from a participating node:
  - ✦ Is constructed as
$$\{T_{ID}, TS, Data\}$$
  - ✦ Where *TS* is the timestamp of local data retrieval. This timestamp is used for synchronizing different occurrences of repeating tasks.

## Examples

- Example:

- ✦ To calculate the total output power (KW) at time  $t_0$  in the entire neighborhood:

- Aggregation plan at node 9 is:

- $\{tid, t_0, power, \{N_5, N_8\}, Enc_K(power) \times I_5 \times I_8, N_{11}, K\}$

- When node 9 receives the aggregation plan:

- 1. It retrieves its own power at  $t_0$ .

- 2. It encrypts the reading with  $K$  to get local input  $C_{p9} = E_K(P_9)$ .

- 3. Node 9 then waits for input from nodes 5,8.

- 4. After receiving  $C_{o5}, C_{o8}$ , node 9 calculates  $C_{o9} = C_{p9} \times C_{o5} \times C_{o8}$ .

- 5. Node 9 submits  $C_{o9}$  to Node 11.

## Analysis

- Comparing:
  - ✦ The in-network aggregation with homomorphic encryption to traditional aggregation approach.
  
- Network:
  - ✦ **Traditional:** messages from all smart meters are routed to collector simultaneously. Let  $\bar{h}$  be the average number of hops for each message to the collector, assuming number of nodes to be  $N$ , total load on the network will be  $\bar{h} * N$ .
  - ✦ **In-network** aggregation, total load will be  $N$  hops.

## Analysis (cont.)

- ○ Scalability, bottleneck and robustness:
  - ✦ Overall scalability highly depends on the smart meter network topology.
  - ✦ In-network aggregation:
    - For a well designed network, the aggregation tree will be wide and shallow. The longest path in an aggregation process is the graph diameter, grows at  $\sqrt{N}$ .
    - Almost **no bottleneck in the in-network aggregation**; since most computations are **distributed**, and also with the **rebalance** scheme.
    - If one start meter fails, **failure is detected immediately** by its parent in aggregation and reported to the collector, the collector **updates** the aggregation tree and **re-issues** the query.

## Analysis (cont.)

- Security and privacy analysis:
  - ✦ The Paillier cryptosystem:
    - **Semantically secure**: polynomial time adversary who intercepts communication **cannot derive significant** information about the plaintext from the ciphers and public key.
    - **Resilient to dictionary attacks**; based on the use of the blinding factor  $r$ , same data will be encrypted to different ciphers with different  $r$ .
    - **WARNING**: all homomorphic encryption systems are malleable; given cipher and public key, an adversary could generate another cipher that decrypts to another meaningful plaintext in the same domain as the original plaintext. Hence, a dishonest meter or fake meter could falsify its data causing inaccurate aggregation result. NOT considered by in-network aggregation, can be solved by increasing physical and software security of smart meters.

## Analysis (cont.)

### ○ Computation:

- ✦ Asymmetric encryption (homomorphic encryption):
  - Is more computationally expensive than symmetric encryption (AES and triple-DES).
  - Traditional (symmetric):
    - Each smart meter encrypt its message, collector to **decrypt N messages**.
  - In-network aggregation:
    - Each smart meter encrypt its message once, and the collector **decrypts one message** (result of aggregation).
    - **Distributes the computation** of the aggregation from collector to intermediate smart meters (with low overhead).

## Personal assessment

- Authors **successfully** extend aggregation concepts from sensor networks into a smart grid framework, carefully handling smart grid issues (smart meters, privacy, etc).
- Authors **fully understand the pros and cons** of their proposed system, and include future research plans to cover the shortcomings.
- Authors **did not provide a quantitative simulation** results that show the gains in savings of computation, and the actual implementation a real smart grid system/subsystem.
- The proposed solution **does not handle the Integrity** aspect in the C-I-A security framework, since authors tried to carefully limit any overhead computations, yet this should be looked at.

## References

1. Homomorphic encryption, [http://en.wikipedia.org/wiki/Homomorphic\\_encryption](http://en.wikipedia.org/wiki/Homomorphic_encryption)
2. What is Homomorphic Encryption, and Why Should I Care? <http://blogs.teamb.com/craigstuntz/2010/03/18/38566/>
3. Fengjun Li; Bo Luo; Peng Liu; , "Secure Information Aggregation for Smart Grids Using Homomorphic Encryption," *Smart Grid Communications (SmartGridComm)*, 2010 First IEEE International Conference on , vol., no., pp. 327-332, 4-6 Oct. 2010 doi: 10.1109/SMARTGRID.2010.5622064 URL: <http://ieeexplore.ieee.org/stamp/stamp.jsp?tp=&arnumber=5622064&isnumber=5621989>