

ECE - Computer Networks II

Problem Set #2

Problem 1. Prove the following properties of the min-plus convolution for one-sided non-decreasing processes F , G , and H :

- (a) *Associativity:* $(F \otimes G) \otimes H = F \otimes (G \otimes H)$.
- (b) *Commutativity:* $F \otimes G = G \otimes F$.
- (c) *Boundedness.* $F \otimes G \leq F$. In particular, $F \otimes F \leq F$.

Solution.

- (a) *Associativity:* The proof can be done for any type of process. We first expand and then rearrange the applications of the ‘inf’ operators, and then perform a substitution $v = u - s$.

$$\begin{aligned}(F \otimes G) \otimes H(t) &= \inf_{u \in \mathbb{R}} \{ \inf_{s \in \mathbb{R}} \{ F(s) + G(u - s) \} + H(t - u) \} \\ &= \inf_{u \in \mathbb{R}} \{ \inf_{s \in \mathbb{R}} \{ F(s) + G(u - s) + H(t - u) \} \} \\ &= \inf_{s \in \mathbb{R}} \{ F(s) + \inf_{u \in \mathbb{R}} \{ G(u - s) + H(t - s - (u - s)) \} \} \\ &= \inf_{s \in \mathbb{R}} \{ F(s) + \inf_{v \in \mathbb{R}} \{ G(v) + H(t - s - v) \} \} \\ &= \inf_{s \in \mathbb{R}} \{ F(s) + G \otimes H(t - s) \} \\ &= F \otimes (G \otimes H)(t)\end{aligned}$$

- (b) *Commutativity:* For any three processes F, G and H , this can be shown by simply performing a substitution $x = t - s$.

$$\begin{aligned}F \otimes G(t) &= \inf_{s \in \mathbb{R}} \{ F(s) + G(t - s) \} \\ &= \inf_{t-x \in \mathbb{R}} \{ F(t - x) + G(x) \} \\ &= \inf_{x \in \mathbb{R}} \{ G(x) + F(t - x) \} \\ &= G \otimes F(t),\end{aligned}$$

where we have used that ‘ $t - x \in \mathbb{R}$ ’ is equivalent to ‘ $x \in \mathbb{R}$ ’.

(c) *Boundedness.* $F \otimes G \leq F$. In particular, $F \otimes F \leq F$.

If G is one-sided, then we have $G \leq \delta$. Then, for a non-decreasing process F , we have $F \otimes G \leq F \otimes \delta = F$. If F is both one-sided and non-decreasing, we can set $G = F$ and get $F \otimes F \leq F$.

Problem 2. Given the functions f and g , compute the convolution $f \otimes g$:

$$f(t) = \begin{cases} 0, & t \leq 1 \\ 2, & 1 < t \leq 2 \\ 2+t, & t > 2 \end{cases}, \quad g(t) = \begin{cases} 0, & t \leq 2 \\ 1+2t & t > 2 \end{cases}$$

Solution.

First, we define $G'(t)$ as time shifted version of $G(t)$:

$$G'(t) = \begin{cases} 0, & t \leq 0 \\ 2t+5 & t > 0 \end{cases}$$

Since we have $G(t) = G' \otimes \delta_2(t)$ by the time shift property of the δ function, the convolution can be rewritten as $F \otimes G(t) = F \otimes G' \otimes \delta_2(t)$. Next, we rewrite $F(t)$ and $G'(t)$ using the indicator function:

$$\begin{aligned} F(t) &= 2 \cdot I_{1 < t \leq 2} + (2+t)I_{t > 2} \\ G'(t) &= (2t+5)I_{t > 0} \end{aligned}$$

First we solve the subproblem of $F \otimes G'(t)$. Later, we apply the time shift $\delta_2(t)$.

$$\begin{aligned} F \otimes G'(t) &= \inf_{0 \leq s \leq t} \{F(s) + G'(t-s)\} \\ &= \inf_{0 \leq s \leq t} \{2 \cdot I_{1 < s \leq 2} + (2+s)I_{s > 2} + (2(t-s)+5)I_{t-s > 0}\} \\ &= \min \left\{ \inf_{0 \leq s \leq 1} \{(2(t-s)+5)I_{s < t}\}, \inf_{1 < s \leq 2} \{2 + (2(t-s)+5)I_{s < t}\}, \right. \\ &\quad \left. \inf_{2 < s \leq t} \{2+s + (2(t-s)+5)I_{s < t}\} \right\} \end{aligned}$$

We next compute the infimums appearing in the above expression:

$$(1) \inf_{0 \leq s \leq 1} \{(2(t-s) + 5) I_{s < t}\}$$

First, observe that if the indicator function $I_{s < t}$ evaluates to one, that is, $s < t$, we have $2(t-s) + 5 > 0$. Therefore, the function $(2(t-s) + 5) I_{s < t}$ achieves the minimum value of zero when $s \geq t$ and the indicator function $I_{s < t}$ evaluates to zero. Since s is in the range $0 \leq s \leq 1$, we can select a value of s that satisfies $s \geq t$ if we have $t < 1$.

In the second case, when $t \geq 1$, the indicator function evaluates to one and we perform the minimization $\inf_{0 \leq s \leq 1} \{2(t-s) + 5\}$. The minimum is obtained for $s = 1$ and is equal to $2t + 3$.

The result of the minimization is:

$$\begin{aligned} \inf_{0 \leq s \leq 1} \{(2(t-s) + 5) I_{s < t}\} &= \begin{cases} 0, & t < 1 \\ 2t + 3, & t \geq 1 \end{cases} \\ &= (2t + 3)I_{t \geq 1} \end{aligned}$$

$$(2) \inf_{1 < s \leq 2} \{2 + (2(t-s) + 5) I_{s < t}\}$$

If the indicator function $I_{s < t}$ evaluates to one, i.e., $s < t$, then $2 + (2(t-s) + 5) > 2$. If the indicator function $I_{s < t}$ evaluates to zero, then the value of the function is equal to 2. Thus, we obtain the minimum of 2 when $s \geq t$, which we can achieve if $t < 2$. In the case when $t \geq 2$ we minimize $\inf_{1 < s \leq 2} \{2 + 2(t-s) + 5\}$. The minimum is obtained for $s = 2$ and is equal to $2t + 3$.

The result of the minimization is:

$$\begin{aligned} \inf_{1 < s \leq 2} \{2 + (2(t-s) + 5) I_{s < t}\} &= \begin{cases} 2, & t < 2 \\ 2t + 3, & t \geq 2 \end{cases} \\ &= \begin{cases} 2, & t < 2 \\ 2 + (2t + 1), & t \geq 2 \end{cases} \\ &= 2 + (2t + 1)I_{t \geq 2} \end{aligned}$$

$$(3) \inf_{2 < s \leq t} \{2 + s + (2(t-s) + 5) I_{s < t}\}$$

If the indicator function $I_{s < t}$ evaluates to zero (for $s = t$) then we have $\inf_{s=t} \{2 + s\} = 2 + t$. If the indicator function $I_{s < t}$ evaluates to one (for $s < t$) then we have $\inf_{2 < s < t} \{2 + s + 2(t-s) + 5\} = \inf_{2 < s < t} \{2t - s + 7\}$, which is minimized for $s \rightarrow t$ and is equal to $t + 7$.

Since $t + 2 < t + 7$, the minimum is obtained for $s = t$ when the indicator function $I_{s < t}$ evaluates to zero. The minimum is

$$\inf_{2 < s \leq t} \{2 + s + (2(t-s) + 5) I_{s < t}\} = t + 2$$

Applying the computed minimums to the convolution we obtain

$$\begin{aligned} f \otimes G'(t) &= \min \{ (2t + 3) I_{t \geq 1}, 2 + (2t + 1) I_{t \geq 2}, t + 2 \} \\ &= \begin{cases} 0, & t < 1 \\ 2, & 1 \leq t < 2 \\ t + 2, & t \geq 2 \end{cases} \end{aligned}$$

Finally we apply the time shift

$$\begin{aligned} F \otimes G(t) &= F \otimes G' \otimes \delta_2(t) \\ &= \begin{cases} 0, & t < 3 \\ 2, & 3 \leq t < 4 \\ t, & t \geq 4 \end{cases} \end{aligned}$$

Problem 3. Define the function $f_n(t) = nI_{t>0}$ for a number $n \in \mathbb{N}$.

1. Show that the function is a traffic process, i.e., f_n is one-sided and non-decreasing.
2. Show the result of the convolution $f_3 \otimes f_2$.

Solution.

A function is one-sided if $f(t) = 0$ for $t \leq 0$. Function $f_n(t) = 0$ for $t \leq 0$ as indicator function $I_{t>0}$ evaluates to zero and thus is one-sided. A function is non-decreasing if $f(t + s) \geq f(t)$ for all $s > 0$.

$$f_n(t + s) = nI_{t+s>0} \geq nI_{t>0} = f_n(t)$$

For $t > 0$ both indicator functions evaluate to 1 (as $s > 0$) and functions are equal. For $t < 0$, if $s \leq -t$ both indicator functions evaluate to 0 making functions equal, and if $s > -t$ first indicator function evaluates to 1 and second to 0 making $f_n(t + s) = n \geq f_n(t) = 0$ as $n \in \mathbb{N}$.

And the convolution is:

$$\begin{aligned} f_3 \otimes f_2(t) &= \inf_{0 \leq s \leq t} \{ f_3(s) + f_2(t - s) \} \\ &= \inf_{0 \leq s \leq t} \{ 3I_{s>0} + 2I_{s<t} \} \\ &= \min \left\{ \inf_{s=0} \{ 2 \}, \inf_{0 < s < t} \{ 3 + 2 \}, \inf_{s=t} \{ 3 \} \right\} \\ &= 2I_{t>0} \end{aligned}$$

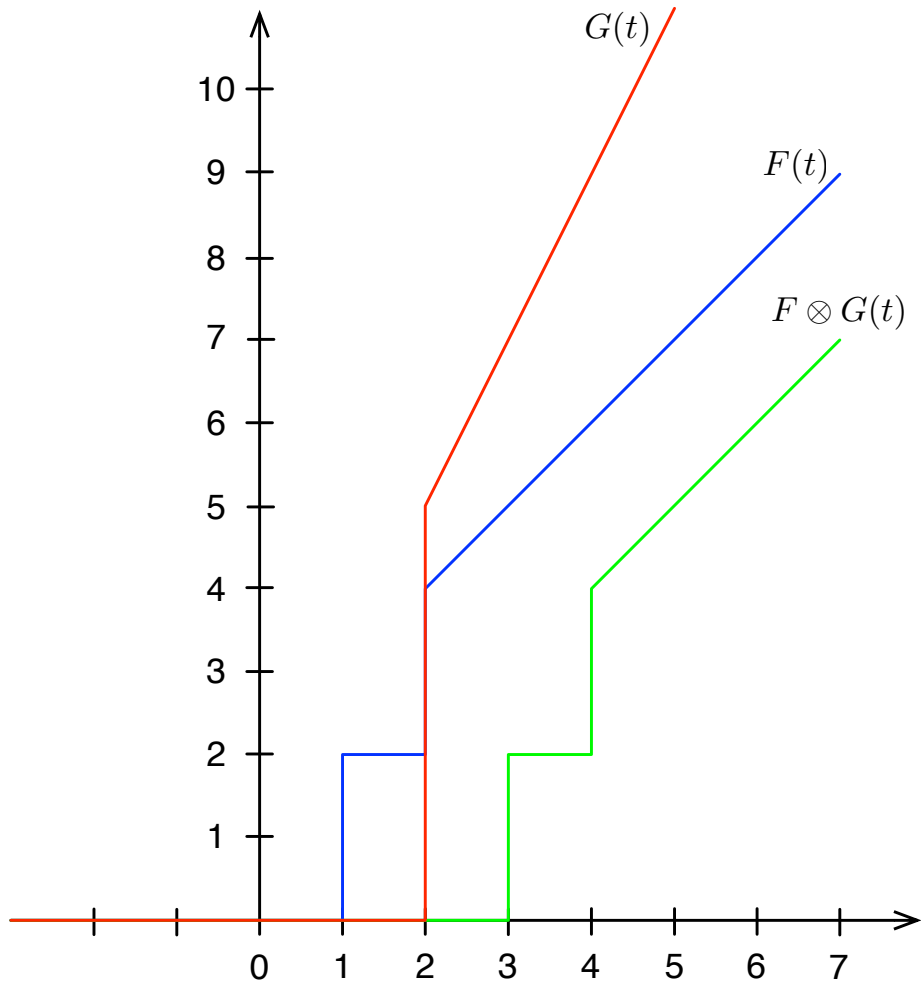


Figure 1: Convolution $F \otimes G(t)$.